



Move	Text Search	Cl	ose
14 SEI L5	96 12:12:15 U.S. Patent & Trac 10 L4 AND TRANSFER(W)REQUEST#	demark Office P000	3
	eit, ab 1-10		
appar	550,997, Aug. 27, 1996, In an interactive itus for preventing inadvertent loading of 7, Tony K. Ip, et al., 395/430; 364/965.76, MAGE AVAILABLE	a programmable read only	
US PA	NO: 5,550,997 [IMAGE AVAILABLE]	L5: 1 of 10	
ROM in ROM in flash deact EPROM predeflash flip flip the f	dand apparatus for loading a ROM image into a dynamic random access memory, what is a second process memory, and in the dynamic random access memory, and in the dynamic random access memory, and in the stored ROM image. A hardward vated by a two-step command procedure so a specifically, a predetermined bit pattern termined address on an address bus, one bit signal to be output to a transistor switch flop. A second predetermined address is reacted to the signal to reach the EPROM. The DM image is loaded into the PROM.	verifying the accuracy of the and receiving a command to be interlock mechanism is as to enable flashing the is written to a first to being latched to cause a hand to be pre-loaded at a and from so as to clock the he transistor switch allowing	y
contr data Krasl	537,626, Jul. 16, 1996, Apparatus for coupol printer operation by transferring controlled printer configuration data between printersky, et al., 395/828; 364/235, 284.4, 93635, 839 [IMAGE AVAILABLE]	ol parameters, printer statu: nter and LAN; Andrew J.	5
US PA	NO: 5,537,626 [IMAGE AVAILABLE]	L5: 2 of 10	
utili: Inter print dispo and a is al for t on the the p	d and apparatus for interfacing a printer the seas a circuit board coupled to the printer face (SCSI) is disposed on the board for the sea and for receiving printer status data for sed on the board, for storing the print day plurality of application programs. A Local so disposed on the board, for receiving the ransmitting the printer status data to the search, for executing the plurality of apprint data to be transmitted to the printer transmitted to the LAN. Preferably, printer intended over the LAN to control printer fundaments.	A Small computer System ransmitting print data to the rom the printer. A RAM is alta, the printer status data, I Area Network (LAN) interface print data from the LAN, at LAN. A processor is disposed plication programs to cause and the printer status data are control data may also be	so ce nd
peripal.,	537,550, Jul. 16, 1996, Interactive networderal statistics with logging level command 395/200.11; 364/264.4, 264.6, 944.9, DIG.1, MAGE AVAILABLEJ	ds; William C. Russell, et	
US PA	NO: 5,537,550 [IMAGE AVAILABLE]	L5: 3 of 10	
inter Netwo the b print	ACT: d and apparatus for logging status information and apparatus for logging status information and a status includes the use of a Small Computer Sypard and having a data channel and a status data to the printer over the data channel, 24 COPY AND CLEAR PAGE, PLEASE	rinter and a Local Area ystems Interface coupled to s channel, for transmitting	5
INP	UT:		

1996 12:12 BERNARD E. GREGORY DE Hold/Res Clr_Out I Ref NDC_Add Pg/Scr_Mode Prt_All Rem Cont_F	Chg_Scr Chg_Scr Chg_Scr
Move Text Search	Close
14 SEP 96 12:12:42 U.S. Patent & Trademark Office	P0004
US PAT NO: 5,537,550 [IMAGE AVAILABLE] requests to the printer and receiving printer status data from the prover the status channel. A memory is coupled to the board and is used storing the received printer status data. A processor is also coupled board for adding beginning and end of print job indicia to the printer prior to transmission to the printer, and for causing the printer state requests to be transmitted to and received from the printer at a first predetermined interval (e.g. every minute). Furthermore, the processor calculates, at a second predetermined interval (e.g. daily) printer statistics based on the received printer status data and the beginning end of job indicia. Finally, the processor stores the calculated printer status statistics in the memory. Preferably, the calculated printer statistics are then stored in a non-volatile memory in the printer at accessed at a later date from a remote location. Also, it is preferably the printer status statistics are provided at a plurality of levels resolution, the levels being selectable from the remote location.	d for d to the data atus st or status ng and nter status nd can be ble if
4. 5,519,695, May 21, 1996, Switch element for fiber channel network Purohit, et al., 370/58.2; 359/117; 370/68.1 [IMAGE AVAILABLE]	ks; Robin
US PAT NO: 5,519,695 [IMAGE AVAILABLE] L5: 4 of 10	
ABSTRACT: A system for facilitating data communications in a fiber channel netroresented. The system comprises a fiber optic switch element which eximplementation of a Fiber Channel network by permitting selective interconnection of a plurality of fiber optic channels. The fiber optic switch element permits both circuit and frame switching. The switch comprises a switch module which is connected to at least one fabric—within the fiber channel fabric. The switch module allocates switching bandwidth. A path allocation system, which is connected to the switch generally directs frames of data between the at least one fabric—port other fabric—ports located within the fabric. A channel module compreport intelligence mechanism is disposed between the switch module and fabric—ports. An element controller provides centralized fabric managements.	nables tic element port ng h module, t to ising a d the
5. 5,513,314, Apr. 30, 1996, Fault tolerant NFS server system and m protocol; David R. Kandasamy, et al., 395/182.04; 364/245.3, 285.3, [IMAGE AVAILABLE]	irroring DIG.1
US PAT NO: 5,513,314 [IMAGE AVAILABLE] L5: 5 of 10	
Anetwork computer system providing for the fault tolerant storage retrieval of data files includes a client system connected to a data communication network that may source a first data transfer request said data communication network for the transfer or retrieval of dat first server system, including first medium for storing data files, connected to the data communication network so as to be responsive to data transfer requests. A second server system, including second med for storing data files is also connected to said data communication to also be responsive to first data transfer requests. A control protocol, established between the first and second server systems, coordinates an asymmetric response by the first and second server sy a first data transfer request, such that file data transferred by the client with the first data transfer request is replicated to the fir and second storing mediums and such that file data transferred to the system in response to the first data transfer is non-replicatively p 12:12:43 COPY AND CLEAR PAGE, PLEASE	to a. A is o first ium network stems to e st e client
INPUT:	

Sep 14,	1996 12:12 BERNARD E. GREGORY THOUSE BERNARD E. GREGORY BERNARD E	Chg_Scr Prt Blk
_ TITCEL T &	Move Text Search Close	
	14 SEP 96 12:13:01 U.S. Patent & Trademark Office P0005	
	US PAT NO: 5,513,314 [IMAGE AVAILABLE] L5: 5 of 10 to the client system by either the first or second server system.	
	6. 5,485,579, Jan. 16, 1996, Multiple facility operating system architecture; David Hitz, et al., 395/200.12; 364/280, 280.6, 280.9, 284, 284.3, 284.4, DIG.1; 395/200.02, 650, 700 [IMAGE AVAILABLE]	
	US PAT NO: 5,485,579 [IMAGE AVAILABLE] L5: 6 of 10	Sign (
	ABSTRACT: This is achieved in a computer system employing a multiple facility operating system architecture. The computer system includes a plurality of processor units for implementing a predetermined set of peer-level facilities wherein each peer-level facility includes a plurality of related functions and a communications bus for interconnecting the processor units. Each of the processor units includes a central processor and the stored program that, upon execution, provides for the implementation of a predetermined peer-level facility of the predetermined set of peer-level facilities, and for performing a multi-tasking interface function. The multi-tasking interface function is responsive to control messages for selecting for execution functions of the predetermined peer-level facility and that is responsive to the predetermined peer-level facility for providing control messages to request or to respond to the performance of functions of another peer-level facility of the computer system. The multi-tasking interface functions of each of the plurality of processor units communicate among one another via the network bus.	
	7. 5,438,528, Aug. 1, 1995, Method and apparatus for testing an interactive network board in a local area network (LAN).; H. Brad Emerson, et al., 364/580, 481, 514B; 370/13, 17; 371/20.1; 395/185.09 [IMAGE AVAILABLE]	
	US PAT NO: 5,438,528 [IMAGE AVAILABLE] L5: 7 of 10	
	ABSTRACT: Method and apparatus for testing an interactive network board having a local area network interface, a Small Computer System Interface, and a test interface comprises supplying power to the interactive board, and performing a power-on self-test program within the interactive board. At the completion of the power-on self-test, a test program is loaded into a RAM on the interactive network board through the test interface, and the test program resident in the RAM is activated. The test program is executed and checkpoint test results are outputted after completion of the test program. A test computer is provided to receive the checkpoint test results and may script additional tests in accordance with checkpoint test results. Preferably, at the completion of the test program, ROM-resident firmware is downloaded into the RAM on the interactive board, and the firmware is loaded from the RAM into a ROM on the interactive network board.	
	8. 5,394,526, Feb. 28, 1995, Data server for transferring selected blocks of remote file to a distributed computer network involving only single data transfer operation; Donald D. Crouse, et al., 395/200.01; 364/225.4, 265.3, 284.4, DIG.1; 395/600 [IMAGE AVAILABLE]	
	US PAT NO: 5,394,526 [IMAGE AVAILABLE] L5: 8 of 10	
	ABSTRACT: A pipelined data server having an improved data transfer architecture is used with a distributed computer network and a plurality of secondary storage devices to efficiently transfer data between the network and the secondary 12:13:01 COPY AND CLEAR PAGE, PLEASE	
	INPUT:	
		4

	1996 12:12 BERNARD E. GREGORY	Chg_Scr
Interrup	Dt Hold/Res Clr_Out Ref NDC_Add Pg/Scr_Mode Prt_All t_Rem Cont_Prt Add_I	
	Move Text Search Clo	se
	14 SEP 96 12:13:18 U.S. Patent & Trademark Office P000	6
	US PAT NO: 5,394,526 [IMAGE AVAILABLE] L5: 8 of 10 storage devices. The pipelined, multiprocessor data server includes a common inter-processor bus that connects one or more communication processors and file processors to one or more device processors, each having a buffer memor as part of the device processor. The common bus provides for global direct access to each of the buffer memories in the device processors by any of the other processors. The buffer memories are also connected to the secondary storage device attached to the device processor by a DMA transfer path in th device processor. In this way, data transfers can occur between the secondar storage device and the network with only one data transfer across the common bus. The data transfer architecture of the device processor of the present invention is also organized in a pipelined manner so as to allow for multiplexed data transfers among the microprocessor, bus interface, I/O controller and buffer memory which comprise the device processor.	y e y
	9. 5,323,393, Jun. 21, 1994, Method and apparatus for obtaining and for controlling the status of a networked peripheral; Lorraine F. Barrett, et al., 370/85.8; 340/825.22 [IMAGE AVAILABLE]	
	US PAT NO: 5,323,393 [IMAGE AVAILABLE] L5: 9 of 10	
	ABSTRACT: Method and apparatus for controlling the same functions of a networked printer that can be manually selected from the front panel of the printer, but remotely through an interactive network board connectable to the printer via a bi-directional printer interface and connectable to a local area network via a local area network interface. A printer status request, for example, from an administrator's console, is issued on the local area networ and directed to the interactive network board to cause the board to interrogate the printer over the printer interface for the status of the manually selectable functions, and to transfer the status of those manually selectable functions from the board onto the local area network. A command t alter the status of those manually selectable functions is issued on the local area network and directed to the interactive board, the command causin the board to transfer the altered status to the printer via the printer interface, whereupon the printer status for the manually selectable function is altered.	k o g
	10. 5,220,501, Jun. 15, 1993, Method and system for remote delivery of retail banking services; Matthew P. Lawlor, et al., 364/408; 379/90; 902/24 [IMAGE AVAILABLE]	
	US PAT NO: 5,220,501 [IMAGE AVAILABLE] L5: 10 of 10	
	ABSTRACT: A practical system and method for the remote distribution of financial services (e.g., home banking and bill-paying) involves distributing portable terminals to a user base. The terminals include a multi-line display, keys "pointer to" lines on the display, and additional keys. Contact is established between the terminals and a central computer operated by a service provider, preferably over a dial-up telephone line and a packet data network. Information exchange between the central computer and the terminal solicits information from the terminal user related to requested financial services (e.g., for billpaying, the user provides payee selection and amount and his bank account PIN number). The central computer then transmits a message over a conventional ATM network debiting the user's bank account in real time, and may pay the specified payees the specified amount electronically or in other ways as appropriate. Payments and transfers may be scheduled in advance or on a periodic basis. Because the central computer 12:13:18 COPY AND CLEAR PAGE, PLEASE	
	INPUT:	

Şep 14,	1996 12:12 BERNARD E. GREGORY THE Hold/Res Clr_Out Image NDC_Add Pg/Scr_Mode Prt_All Emt_Rem Cont_Prt Add_Blk Image Rem Cont_Prt Add_Blk Image Rem Cont_Prt Add_Blk Image Rem R	Chg_S Prt B
III dell'il	Move Text Search Close	
	14 SEP 96 12:13:34 U.S. Patent & Trademark Office P0007	
	US PAT NO: 5,220,501 [IMAGE AVAILABLE] L5: 10 of 10 interacts with the user's bank as a standard POS or ATM network node, no significant software changes are required at the banks' computers. The terminal interface is extremely user-friendly and incorporates some features of standard ATM user interfaces so as to reduce new user anxiety.	
	=> file jpoabs FILE 'JPOABS' ENTERED AT 12:13:42 ON 14 SEP 96	

	* * CURRENTLY, DATA IS LOADED THROUGH THE ABSTRACT PUBLICATION * * DATE OF DECEMBER 26, 1994 * * THE LATEST GROUPS RECEIVED ARE: C1292 E1651, M1731 & P1851. * * * * * * * * * * * * * * * * * * *	
	=> dis his	
	(FILE 'USPAT' ENTERED AT 12:08:56 ON 14 SEP 96)	
	SET PAGELENGTH 62 SET LINELENGTH 78 L1 101 S HYPERTEXT OR HYPER(W) TEXT OR HYPERLINK### OR HYPER(W) LINK## L2 356 S HTTP OR INTERNET OR INTER(W) NET L3 454 S L1 OR L2 L4 405 S L3 AND COMPUTER# L5 10 S L4 AND TRANSFER(W) REQUEST#	
	FILE 'JPOABS' ENTERED AT 12:13:42 ON 14 SEP 96	
	=> s l1	
	1 HYPERLINK### 263 HYPER 50189 LINK### 1 HYPER(W)LINK### 67 HYPERIEXT OR HYPER(W)TEXT OR HYPERLINK### OR HYPER(W)LINK###	
Section 1	=> s 12 0 HTTP 12 INTERNET 14055 INTER 13287 NET	
	7 INTER(W)NET L7 19 HTTP OR INTERNET OR INTER(W)NET	
	=> s 16 or 17 L8 86 L6 OR L7	
	=> s 18 and computer# 61465 COMPUTER#	
	L9 4 L8 AND COMPUTER# 12:14:23 COPY AND CLEAR PAGE, PLEASE	
	INPUT:	

	RNARD E. GREGO		Mode Prt_All	Rem[Cont Prt]Ad	Cha Cha
Move Move		Text Sea			Close
14 SEP 96 12	2: 14: 31	U.S. Patent	& Trademark Offi	ce P(8000
=> d cit,ab					
1. 07-22579	95, Aug. 22 , 19	95, KNOWLEDGE AC	CUMULATING AND S G06F 17/30	HARING DEVICE;	
07–225795			L9:	1 of 4	
ABSTRACT:					
the members knowledge is	of a community accumulated, regional commu	can easily view regarding a know	ing and sharing multimedia info ledge accumulati he information m	rmation in which ng and sharing	1
composed of control in a by the hyper information etc., by down device 5, an image assist characterist	a host station regional communitation structure composing the particular and interface againg the retries of a member	a 1 performing the nunity, etc., and a based on multime part and a view or ogram from this ent 9 provided with action of the state of t	accumulating and the knowledge accumulating the legislation of the legislation accumulation and the learning of the learning and accumulation and accumulation ac	mulation and program 8 compose or the ing the program In the viewing unction of a move individual	sed 8, Ving
COPYRIGHT:	(C) 1995, JP0				
2. 07-10495 B41J 5/30	57, Apr. 21, 19	995, PRINTER SERV	ER; SHOJI KATSUR	AGAWA, G06F 3/12	2;
07-104957			L9:	2 of 4	
ABSTRACT:					
provided wit		erver by process	onstruction of a ng the packets o		nds
internet 1 in a recepting discriming packet is treather the protocol of each protocol	is received by on buffer 5. I hated by a protest to construct to color and representations.	a data receiving The protocol of to Tocol discriminations of protocol p t processing part Toduce print data	a host computer, part 3, and tem the packet of the ing part 6, and processing parts as 7-9 extract date. Then, the print part 10 to an	porarily stored reception buffer the discriminate 7 -9 according to the pack to data are	er 5 ed to cet
COPYRIGHT:	(C) 1995, JP0				
3. 05-21660 HYPER; MICHI	07, Aug. 27, 19 KO OKAMOTO, GO	993, SYSTEM FOR (96F 3/14; //G06F	ENERATING TABLE	OF CONTENTS FOR	
05-216607			L9:	3 of 4	
12:14:38 COP	PY AND CLEAR PA	AGE, PLEASE			
INPUT:					
Si granda de la companya de la comp					

	996 12:14 BERNARD E. GRE		All int Bon Cont But		g_Scr
Interru	Move	ef NDC_Add Pg/Scr_Mode Prt_ Text Search	ATT C_Rem Cont_FFT	Close	t_bik
			1 000:		
	14 SEP 96 12:14:57	U.S. Patent & Tradema		P0009	
	05–216607		L9: 3 of 4		
	ABSTRACT:				
	PURPOSE: To provide the processor handling cards of a computer.	means displaying any card for displaying the graphical i	or the hyper text nformation on the so	creen	
	card management section prevent buttons from be- button information sect- procedure moving to the holds the image of a fig to the center coordinate	d information section 202 is of card number N to calcula ing overlapped each other on ion 207 is prepared to make (i)th card. A display image gure (i) to calculate the spe of a button information se keeps the center coordinate is added to card set S.	the the space so as to the section 202. A a procedure section information section acceptance of the button acceptance 207. A position	to new 209 a 1 211 djacent on	
	4. 02-187819, Jul. 24, SUZUKI, et al., G06F 3/2	1990, DISPLAY SEQUENCE CONT	ROL SYSTEM; TOSHIMI	rsu	
	02-187819		L9: 4 of 4		
	ABSTRACT:				
an, maliji	that no time series info	ime loss to refer doubly by primation of a link destination exists in the same time somation of the link destination.	on is displayed doub series and also, the	oly part	
	device 2, and also, an econsisting of plural pice computer 4. At this time of the same time series the plural pieces of time middle way of the sequence part of the time series link origin in point of displayed on the displayinformation of the link access by the input devices	external memory device 3 which external memory device is used information is comprised of the series information, and and external exte	ch stores the part tion are controlled a sed, and the sequence the part consisting also, a link exists information, and who stination exists behind, the sequence is part of the time ser a the device 1 by male colled to that the part	oy a e g of on the en the ind a ies cing	
	=> dis his			,	
	(FILE 'USPAT' ENTE	RED AT 12:08:56 ON 14 SEP 96	5)		
	SET LINI L1 101 S HYPER: L2 356 S HTTP (L3 454 S L1 OR L4 405 S L3 ANI	O COMPUTER# O TRANSFER(W)REQUEST#	ERLINK### OR HYPER(W	LINK##	
	INPUT:		WELVE OF THE STREET		
				- [

Serial Serial Close		1996 12:15 BERNARD E. GREGORY Hold/Res Clr_Out Ref NDC_Add Pg/Scr_Mode Prt_All t_Rem Cont_Prt	Add_Blk
FILE 'JPOABS' ENTERED AT 12:13:42 ON 14 SEP 96 L6 67 S L1 L7 19 S L2 L8 86 S L6 OR L7 L9 4 S L8 NOT COMPUTER# => \$ 18 not 19 L10 82 L8 NOT L9 => \$ 110 and transfer (w) request# 142790 TRANSFER 41715 REQUEST# 1473 TRANSFER (W) REQUEST# L11 0 L10 AND TRANSFER (W) REQUEST# L12 8 L10 AND (TRANSFER OR REQUEST#) => \$ 110 and (transfer or request#) 142790 TRANSFER 41715 REQUEST# 2 \$ L10 AND (TRANSFER OR REQUEST#) => d cit, ab 1-8 1. 06-338855 Dec. 6, 1994, [NTHENNIM], TERMINATION SYSTEM, INTERMEDIATE SYSTEM AND [NTHENNIM] COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06F 13/00; H04L 12/36 06-338855 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmissing is greened from a host layer flows hown in a figure), a data bucket transmission control part 13 retrieves the routing table, and the information held beyond the holding time is time for letting Excess is made from a host layer flows hown in a figure), a data bucket transmission control part 13 retrieves the routing table, and the information system is present. When it is registered, the corresponding data link layer is read, and the data bucket transmitting Excess is made from a host layer flows the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket transmited to the addre		Move Text Search	Close
L6 67 S L1 L7 19 S L2 L8 86 S L6 OR L7 L9 4 S L8 NOT L9 => S 18 not 19 L10 82 L8 NOT L9 => S 110 and transfer (w) request# 142790 TRANSFER 41715 REQUEST# 1473 TRANSFER (w) REQUEST# 1473 TRANSFER (w) REQUEST# 1473 TRANSFER (w) REQUEST# => S 110 and (transfer or request#) 142790 TRANSFER 41715 REQUEST# L12 8 L10 AND (TRANSFER OR REQUEST#) => d cit, ab 1-8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting promises is made from a lost layer (not shown in a figure) a data bucket transmitting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting promises the Name of the address of the address destination termination system is present. When it is registered the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681		14 SEP 96 12:15:24 U.S. Patent & Trademark Office	P0010
L7 L9 L8 86 S L6 OR L7 L9 4 S L8 AND COMPUTER# >> s 18 not 19 L10 82 L8 NOT L9 >> s 110 and transfer(w)request# 142790 TRANSFER 41715 REQUEST# 1473 TRANSFER (W)REQUEST# 1473 TRANSFER (W)REQUEST# 25 s 110 and (transfer or request#) 142790 TRANSFER 41715 REQUEST# L11 2		FILE 'JPOABS' ENTERED AT 12:13:42 ON 14 SEP 96	
L10 82 L8 NOT L9 >> \$ 110 and transfer (w) request# 142790 TRANSFER 14715 REQUEST# 1473 TRANSFER TWO REQUEST# L11 0 L10 AND TRANSFER (W) REQUEST# >> \$ 110 and (transfer or request#) 142790 TRANSFER 14715 REQUEST# L12 8 L10 AND (TRANSFER OR REQUEST#) >> d cit,ab.1-8 1. 06-338855, Dec. 6, 1994, INTERNET, TERMINATION SYSTEM, INTERMEDIATE SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06F 13/00; H04L 12/36 06-338855 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting requires its made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the ad		L7 19 S L2 L8 86 S L6 OR L7	
14790 TRANSFER 41715 REQUEST# 1473 TRANSFER (WREQUEST# 1473 TRANSFER (WREQUEST# 2> s 110 and (transfer or request#) 142790 TRANSFER 41715 REQUEST# L12 8 L10 AND (TRANSFER OR REQUEST#) 2> d cit,ab 1-8 1. 06-338885. Dec. 6, 1994, INTERNET, TERMINATION SYSTEM, INTERMEDIATE SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06F 13/00; H04L 12/36 06-338885 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmissions whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681 L12: 2 of 8 12:16:18 COPY AND CLEAR PAGE, PLEASE	n en		
142790 TRANSFER 41715 REQUEST# L12 8 L10 AND (TRANSFER OR REQUEST#) => d cit, ab 1-8 1. 06-338885, Dec. 6, 1994, INTERNET, TERMINATION SYSTEM, INTERMEDIATE SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06F 13/00; H04L 12/56 06-338885 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting INDURST is made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or i		142790 TRANSFER 41715 REQUEST# 1473 TRANSFER(W)REQUEST#	
1. 06-338885, Dec. 6, 1994, INTERNET, TERMINATION SYSTEM, INTERMEDIATE SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06F 13/00; H04L 12/56 06-338885 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting Sequest is made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681 L12: 2 of 8 12:16:18 COPY AND CLEAR PAGE, PLEASE		142790 TRANSFER 41715 REQUEST#	
SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, HO4L 12/28; G06F 13/00; H04L 12/56 06-338885 L12: 1 of 8 ABSTRACT: PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting request is made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system or intermediate system of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system is present. When it is a system is not a system is present.		=> d cit,ab 1-8	
PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681 L12: 2 of 8 12:16:18 COPY AND CLEAR PAGE, PLEASE		SYSTEM AND INTERNET COMMUNICATING METHOD; MASASHI KUDO, H04L 12/28; G06	3
PURPOSE: To exchange routing information without any excess load by stopping the multicast of a connection information bucket from an intermediate system to each termination system, and inquiring the address of the data link layer of the destination to the termination system in which the necessity of transmission is generated. CONSTITUTION: A routing table 12 stores routing information, that is, the N layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting request is made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681 L12: 2 of 8 12:16:18 COPY AND CLEAR PAGE, PLEASE		06-338885 L12: 1 of 8	
layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it, and holding time. The holding time is time for letting the routing information hold in the routing table, and the information held beyond the holding time is deleted. When a data bucket transmitting request is made from a host layer (not shown in a figure), a data bucket transmission control part 13 retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When it is registered, the corresponding data link layer is read, and the data bucket is transmitted to the address destination termination system or intermediate system of the address. COPYRIGHT: (C)1994, JPO 2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681 L12: 2 of 8 12:16:18 COPY AND CLEAR PAGE, PLEASE		PURPOSE: To exchange routing information without any excess load by stathe multicast of a connection information bucket from an intermediate sto each termination system, and inquiring the address of the data link of the destination to the termination system in which the necessity of	ystem
2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., H04L 29/06; G06F 13/00; H04L 12/28 06-152681		layer address of a destination termination system, the data link layer address of the direct transmission destination node corresponding to it holding time. The holding time is time for letting the routing informat hold in the routing table, and the information held beyond the holding is deleted. When a data bucket transmitting request is made from a host layer (not shown in a figure), a data bucket transmission control part retrieves the routing table 12, and confirms whether or not the N layer address of the address destination termination system is present. When registered, the corresponding data link layer is read, and the data bucket transmitted to the address destination termination system or intermedian	, and ion time 13 it is cet is
29/06; G06F 13/00; H04L 12/28 06-152681		COPYRIGHT: (C)1994, JPO	
12:16:18 COPY AND CLEAR PAGE, PLEASE		2. 06-152681, May 31, 1994, CMIP-SNMP GATEWAY; TAKASHI KAGEI, et al., 1 29/06; G06F 13/00; H04L 12/28	H04L
		06-152681 L12: 2 of 8	
INPUT:		12:16:18 COPY AND CLEAR PAGE, PLEASE	
		INPUT:	
			-

Intérru	1996 12:16 BERNARD E. GREGORY Best Be	nt_Prt Add_Blk
	Move Text Search	Close
	14 SEP 96 12:16:31 U.S. Patent & Trademark Office	P0011
	06-152681 L12: 2 of	8
	ABSTRACŢ:	
. A	PURPOSE: To manage an inter-net with OSI management by providing SNMP gateway between an OSI manager and an SNMP agent and applyin processing to the management information specified by an inter-ne management into management information regulated by the OSI manag	g conversion
	CONSTITUTION: A CMIP-SNMP gateway 1 exchanges management informat OSI manager 2 via an OSI communication interface 22 and exchanges management information with an SNMP agent 3 via an ICP/IP interfa A-associate processing 12, an A-release processing 13, an A-abort 14, an m-get processing 15, an m-set processing 16, and an m-even processing 17 implement the conversion of CMIP protocol data into expression used for a gateway 1 and conversion of the internal ex into CHIP protocol data. A get-request processing 18, a get-respo processing 19, a set-request processing 20 and a trap processing implement the conversion of SNMP protocol data to the internal ex the conversion of the internal expression into SNP protocol data.	the ce 23. An processing t report an internal pression nse 21 pression and
	3. 06-120944, Apr. 28, 1994, OPTIMUM MESSAGE TRANSFER ROUTE DECI SYSTEM IN INTERNET ENVIRONMENT; SHINJI MATSUMOTO, HO4L 12/00; G06	DING F 13/00
	06-120944 L12: 3 of	8
	ABSTRACT:	
	PURPOSE:To accurately select a transfer route with shortest dela an internet environment.	y time in
	constitution: A delay time information table 1-4 stores transfer with an adjacent gateway and an optimum transfer route table 1-5 optimum transfer route to the adjacent gateway. A delay time info exchange part 1-1 transmits/receives messages for measuring the d and delay time information table messages with the adjacent gatew predetermined time interval and a delay time information analysis calculates the delay time with the adjacent gateway from the recemessages for measuring the delay time and updates the delay time table. A transfer route decision part 1-3 decides the optimum transfer route to the adjacent gateway from the received delay time information table its own transfer route table and its own delay time information table	stores an rmation elay time ay at a part 1-2 ived information nsfer ation table and updates
	4. 06-110926, Apr. 22, 1994, INFORMATION RETRIEVING DEVICE; SHIG G06F 15/40; G06F 1 2/00	ERU SUZUKI,
	06-110926 L12: 4 of	8
	ABSTRACT:	
	PURPOSE: To speedily display the information of the next node whe node is moved to the next node.	n a certain
	CONSTITUTION: In a hyper text system, when a display request is i to a certain node, a reading means 11 reads the node. A link rela 12:16:31 COPY AND CLEAR PAGE, PLEASE	ssued tion
	INPUT:	
		

_Scr _Blk

1996 12:16 BERNARD E. GREGORY The Hold/Res Clr_Out Ref NDC_Add Pg/Scr_Mode Prt_Al Ref Ref NDC_Add Pg/Scr_Mode Prt_Al Ref Ref NDC_Add Pg/Scr_Mode Prt_Al Pg/Scr_Mode Pt/Scr_Mode Pt/	em[Cont_Prt[Add B]k	Chg_Scr
Move Text Search	Close	
14 SEP 96 12:16:48 U.S. Patent & Trademark Office	P0012	
	4 of 8	
discriminating means 12 discriminates a link relation from the information described in the read information. When the link certain node is discriminated, a first reading means 13 reads of the next node, and constitutes the screen image of the not the instruction of the next node.	relation from a sthe information	
 5. 04-326443, Nov. 16, 1992, INTERFACE DEVICE; NORIO TAKAGI, G06F 13/10	, G06F 13/00;	
04-326443 L12: 5	5 of 8	
ABSTRACT:		
PURPOSE:To obtain an interface device which accesses a modul TCP/IP as an external input/output circuit (external I/O) and TCP/IP to a program generated in an operating system where the device for supplying the communication service of TCP/IP is a	d which supplies he interface	
CONSTITUTION: A command input part 121 inputs a command from program and a command analysis part 122 analyzes the command service request part 123 generates a control protocol/interneservice request corresponding to the analyzed result to a TCI and a service response from the TCP/IP module 20 corresponding request is inputted in a TCP/IP service response input part corresponding to the service response is outputted from a respont 125 to the application program.	. A TCP /IP T protocol P/IP module 20, ng to the service 124. A response	
6. 04-167841, Jun. 15, 1992, LAN SYSTEM; YUTAKA HIRASAWA, HO	04L 12/40	
04-167841 L12: 6	6 of 8	,
ABSTRACT:		
PURPOSE: To improve the communication processing capability of side up to its own limit by connecting plural LAN interface of device so as to decentralize a load.	of an equipment devices to one	
CONSTITUTION: An ARP request received by a LAN controller LCI the reception of a multiple address frame is allowed is given processing section 24. The processing section 24 checks whether the processing section 24 controlled the LCP when coincident, a load table 22 is referenced and the LCP when coincident, a load table 22 is referenced and the LCP when selected as an ARP request reply source among LCP30-1-30-1 LCP30-n). The processing section 24 generates the ARP reply address of the LCP 30-n as a sender address. A reply of sendance and the LCP30-n via a land cable 11 is requested to the LCP30-n via	n to an ARP her or not an own other hand, ith a least load n (e.g. taking the LAN ing the generated	
7. 03-149634, Jun. 26, 1991, HYPERTEXT SYSTEM; TETSUO MURANA G06F 12/00; G06F 9/44	AGA, et al.,	
03-149634 L12: 1	7 of 8	
12:16:48 COPY AND CLEAR PAGE, PLEASE		
		-
INPUT:		

	996 12:16 BERNARD E.		Mada Dat All Dat Dan Con		Scr
Interru	Move	_Ref NDC_Add Pg/Scr Text Se		t_Prt Add_Blk Prt	_DIK

	14 SEP 96 12:17:05	U.S. Patent	: & Trademark Office	P0013	
	03–149634		L12: 7 of 8		
	ABSTRACT:				
	hypertext by automa-	possible for a reader tically selecting and th the reader's reques	to efficiently read out displaying a node to be read or knowledge.	a ead out	
	determine a node to be observed by the a to be displayed nex- access storage part In the case of read line applied by an determined by the sy knowledge even if the	be observed next with reader, a hypertext control based upon default to and displays the defined out the hypertext author, the node to be ystem side in matching he reader forcedly sel	instruction for allowing a nout specifying the node report of the node of the note of the node of the note of the note of the node of the	required to the node a node olay part 3. a guide tically or the node	
	8. 02-224431, Sep. TANEDA, H04L 12/28;	6, 1990, ROUTING INFO H04L 12/56	ORMATION CONTROL SYSTEM; K	CAZUMASA	
	02-224431		L12: 8 of 8	3	
	ABSTRACT:				
	improve the reliabi constitution accord received packet from	lity of the whole IN b ing to the sending ori m a distant network to	I the flexibility of IN, a by detecting the change of gin network address infor which a network layer re dating routing information	internet rmation of a epeating	
	processing, it information. The rout. IN constitution, and to update the routing issues a routing propert 1 executes routing the routing informing the transport of the routing information in the routing in the routing information in the routing information in the routing information in the routing in the routing information in the routing in the routing in the routing in the routing in the rou	rms a routing informating information control d after it executes ro ng information in a ro ocessing request to a ting processing based rmation storage part 4 mitting receiving part	ving part 2 executes reception control part 3 of pace of part 3 detects the charmonic part information control puting information storage routing part 1. After the on the updated routing in 1, it issues a sending part 2 of destination PA and as, the automatically added to	cket age of the aprocessing a part 4, it a routing aformation alest by LID	
	=> file uspat FILE 'USPAT' ENTERE	D AT 12:17:19 ON 14 SE	EP 96		
	* U.S.	* * * * * * * * * * * * * * * * * * *	FILE *		
	=>				
	TAIDLIT .		***		
	INPUT:				
**************************************					siiii00000000